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Logic

Antidirected paths in 5-chromatic digraphs

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Abstract

Let T_5 be the regular 5-tournament. B. Grünbaum proved that T_5 is the only 5-tournament which contains no copy of the antidirected path P_4 . In this Note, we prove that, except for T_5 , any connected 5-chromatic oriented digraph in which each vertex has out-degree at least two contains a copy of P_4 . It will be shown, by an example, that the condition that each vertex has out-degree at least two is indispensable. *To cite this article: A. El Sahili, C. R. Acad. Sci. Paris, Ser. I 339 (2004).* © 2004 Académie des sciences. Published by Elsevier SAS. All rights reserved.

Résumé

Chemins antidirigés dans les graphes 5-chromatiques. Soit T_5 le tournoi régulier contenant cinq sommets. B. Grünbaum a prouvé que T_5 est le seul 5-tournoi qui ne contient pas le chemin antidirigé P_4 . Nous prouvons dans cette Note que T_5 est le seul graphe orienté 5-chromatique dans lequel tout sommet a un degré extérieur au moins deux qui ne contient pas le chemin antidirigé P_4 . On prouve à l'aide d'un exemple que la condition « tout sommet a un degré exterieur au moins deux » est indispensable. *Pour citer cet article : A. El Sahili, C. R. Acad. Sci. Paris, Ser. I 339 (2004).* © 2004 Académie des sciences. Published by Elsevier SAS. All rights reserved.

1. Introduction

The digraphs considered here have no loops or multiple edges. An *oriented graph* is a digraph in which, for every two vertices x and y, at most one of (x, y), (y, x) is an edge. The digraphs used in this Note are all oriented graphs. By G(D) we denote the underlying graph of a digraph D. The chromatic number of a digraph is the chromatic number of its underlying graph. A graph G is k-critical if $\chi(G) = k$ and $\chi(G - v) = k - 1$ for any vertex v in V(G).

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A block of an oriented path is a maximal directed subpath. We recall that the *length* of a path is the number of its edges. The antidirected path is an oriented path in which each block is of length 1. We denote by P_n the antidirected path of length n, beginning with a backward edge.

The problem of determining which oriented paths lie in a given *n*-chromatic digraph *D* is a well-known problem. When *D* is an *n*-tournament, the problem has been completely resolved (Havet and Thomassé [6]). However, the case of an arbitrary *n*-chromatic digraph is still an open question. We know only that an *n*-chromatic digraph contains a directed path of length n - 1 (Roy [7], Gallai [4]), and a path of length n - 1 formed by two blocks, one of which has length 1 [2]. In this Note, we will be interested in the antidirected paths. In order to generalize the results found on tournaments to arbitrary digraphs, and as a first step in this direction, we generalize to 5-chromatic digraphs a particular result of Grünbaum on 5-tournaments: any 5-tournament, except for the regular tournament T_5 , contains a copy of P_4 .

2. The main result

Theorem 2.1. Let *D* be a 5-chromatic connected digraph, distinct from T_5 , in which each vertex has out-degree at least two. Then *D* contains a copy of P_4 .

To prove this theorem, we need several lemmas.

Lemma 2.2 (Grünbaum [5]). Except for T₅, any 5-tournament contains a copy of P₄.

Corollary 2.3. Let D be as in the above theorem. If D contains T_5 , then D contains a copy of P_4 .

In the sequel, D will denote an oriented digraph as described in Theorem 2.1; by the above corollary we may assume that D contains no 5-tournament as a subdigraph. Moreover, we suppose to the contrary that D contains no copy of P_4 . Let D' be a 5-critical subdigraph of D and let D° be the subdigraph of D' induced by the vertices of out-degree at least three in D'.

Let *G* be a graph which contains no K_{2n+1} , where $n \ge 2$. Suppose that we can orient *G* in such a way that each vertex has in-degree at most *n*. It is shown in [1] that $\chi(G) \le 2n$. We have then the following lemma

Lemma 2.4. The set $V(D^{\circ})$ is not empty.

Lemma 2.5. Let v be a vertex of D and let x, y be two vertices in $N^-(v)$. If $x \in V(D^\circ)$, then $y \notin V(D^\circ)$.

Corollary 2.6. For every vertex v in D° , $d_{D^\circ}^-(v) \leq 1$.

Lemma 2.7. Let *H* be a connected digraph in which each vertex has in-degree at most one. Then *H* contains at most one cycle.

Lemma 2.8. Let v be a vertex of D such that $d^+(v) \ge 3$ and let x, y and z be three distinct vertices in $N^+(v)$. Suppose that $x \to y$. Then $N^-(y) = N^-(z) = \{v, x\}$.

We may easily deduce that $x \to z$ and $yz \notin E(G(D))$ in this case.

Corollary 2.9. Let x and y be two adjacent vertices of D. Suppose that there exist two vertices v and v' of D such that $\{x, y\} \subseteq N^+(v) \cap N^-(v')$. Then $N^+(v) = \{x, y\}$.

Lemma 2.10. The set $V(D^{\circ})$ is independent in D.

Claim 1. Any connected component L of D° contains a vertex v such that $N^+(v) \cap (V(D') \setminus V(D^{\circ}))$ has at least two vertices.

Proof. If *L* is a cycle, then each vertex of *L* satisfies the claim; otherwise *L* contains a vertex *v* of out-degree zero in D° , and so $N^+(v) \subseteq V(D') \setminus V(D^\circ)$. \Box

Proof of Lemma 2.10. Suppose to the contrary that D° is not an independent set, then there is a connected component L of D° containing at least two vertices. We can choose a vertex v in L satisfying the claim such that $d_L^-(v) = 1$. Let v' be a vertex in L such that $v' \to v$ and let v_1 , v_2 and v_3 be three vertices in $N_{D'}^+(v)$ such that $\{v_1, v_2\} \subseteq V(D') \setminus V(D^{\circ})$. The digraph D' is 5-critical, so any vertex has degree at least 4 in D'. Since for any $i \in \{1, 2\}, d_{D'}^+(v_i) \leq 2$, we have $d_{D'}^-(v_i) \geq 2$. Therefore, there is a vertex u of D' and $j \in \{1, 2\}$ such that $u \notin \{v, v_1, v_2\}$ and $u \to v_j$; we have either $u \notin \{v, v_1, v_2, v_3\}$ or $u = v_3$. In the latter case $v_3 \notin V(D^{\circ})$ by Lemma 2.5. We have $d_{D'}^-(v_3) \geq 2$, so there is a vertex w of D' such that $w \notin \{v, v_1, v_2, v_3\}$ and $w \to v_3$, thus we may assert that there exists a vertex u of D' and $j \in \{1, 2, 3\}$ such that $u \notin \{v, v_1, v_2, v_3\}$, $v_j \notin D^{\circ}$ and $u \to v_j$. Let u' be a vertex of D distinct from v_j such that $u \to u'$. If $u' \neq v$, the path $u'uv_jvv_h$ is a copy of P_4 , where $h \in \{1, 2, 3\} \setminus \{j\}$ is chosen such that $u' \neq v_h$, a contradiction. Otherwise, let w be a vertex in $N^+(v') \setminus \{v, v_j, u\}$. Such a vertex exists since $d^+(v') \geq 3$ and $v_j \notin N^+(v')$ by Lemma 2.5. The path $v_juvv'w$ is a copy of P_4 , a contradiction. \Box

In the sequel, we will need the following theorem proved by Gallai [3].

Theorem 2.11. Let G be a k-critical graph, where k is a positive integer. Let G_m be the subgraph of G induced by the vertices of degree k - 1. Then each block of G_m is either complete or a chordless odd cycle.

 D_4 will denote the subdigraph of D' induced by the vertices of degree 4.

Lemma 2.12. Any vertex of D' has in-degree (in D') at least 2.

We now associate to each vertex v in D° the set

$$S(v) = \{t(v), t'(v), v_0, \dots, v_{g(v)}, v_{g(v)+1}\}, \quad 0 \le g(v) \le 5,$$

defined as follows (see Fig. 1): $\{v_0, t(v), t'(v)\} = N_{D'}^+(v)$ where $v_0 \to t(v)$ and $v_0 \to t'(v)$, $v_1 = v$. Set $T(v) = \{t(v), t'(v)\}$. If $d_{D'}^-(v_0) \ge 3$, put g(v) = 0; if not, let v_2 be the unique vertex of D' distinct from v_1 such that $v_2 \to v_0$. We have $v_2 \to v_1$. Again, if $d_{D'}^-(v_1) \ge 3$, put g(v) = 1; otherwise, let v_3 be the unique vertex of D' distinct from v_2 such that $v_3 \to v_1$; such a vertex exists by Lemma 2.12. We have $v_3 \to v_2$, since otherwise we have either a path P_4 in D or $d_{D'}^-(v_0) \ge 3$. We may continue this process until meeting the first vertex of in-degree at least three in D'; call this vertex $v_{g(v)}$, where g(v) is the number of iterations required. Such a vertex exists and $g(v) \le 5$. In fact, suppose that v_1, \ldots, v_5 are defined as above and $d_{D'}^+(v_i) = 2$, $i = 1, \ldots, 4$. By Corollary 2.9,



Fig. 1. The case g(v) = 5.

we have $d_{D'}^+(v_i) = 2$, i = 2, ..., 5. If $d_{D'}^-(v_5) = 2$ the vertices $v_2, ..., v_5$ will be in the same block of D_4 . By Theorem 2.11, $D'[v_2, ..., v_5]$ is complete, which is a contradiction since $v_2v_5 \notin E(G(D))$.

Set $O(v) = t\{z \in D': z \neq v_{g(v)+1} \text{ and } z \to v_{g(v)}\}$; we have $z \to v_{g(v)+1}$ for every z in O(v).

Lemma 2.13. Let u and v be two distinct vertices of D° . We have:

 $S(u) \cap S(v) = \phi.$

Lemma 2.14. Set $L = \{v_{g(v)}: v \in D^{\circ}\}$. We have:

(i) $d_{D'}^{-}(x) = 3$ for any x in L.

(ii) $d_{D'}(x) = 2$ otherwise.

Corollary 2.15. For any vertex v in D° , O(v) contains exactly two vertices.

Proof of Theorem 2.1. Define the sets:

$$S = \bigcup_{v \in V(D^\circ)} S(v), \quad O = \bigcup_{v \in V(D^\circ)} O(v), \quad T = \bigcup_{v \in V(D^\circ)} T(v).$$

We have $|O| \leq |T|$. If O = T, then $N_{D'}(v) \subseteq S$ for every v in S. Since D' is critical, it must be connected and so D' = D'[S]. We define a colouring c of D' as follows: Let v be a vertex in D° . Put c(t(v)) = c(t'(v)) = 1, $c(v_0) = 2$, $c(v_1) = 3$. If g(v) = 1, put $c(v_2) = 4$. If g(v) > 1, the colours 1, 2 and 3 suffice to colour $S(v) \setminus \{v_{g(v)}, v_{g(v)+1}\}$. Put $c(v_{g(v)}) = 4$ and $c(v_{g(v)+1}) = i$ where $i \in \{2, 3\}$ is chosen such that $i \neq c(v_{g(v)-1})$. It is clear that c is a proper 4-colouring of the 5-chromatic digraph D', a contradiction.

If $O \neq T$ then, since $|O| \leq |T|$, there is a vertex v in D° such that either $t(v) \notin O$ or $t'(v) \notin O$. Suppose, without loss of generality, that $t(v) \notin O$. Then $N_{D'}^+(t(v)) \cap S = \phi$. Let $N_{D'}^+(t(v)) = \{u, u'\}$. We have $\{u, u'\} \cap (D^{\circ} \cup L) = \phi$, so $d_{D'}^-(u) = d_{D'}^+(u) = d_{D'}^-(u') = d_{D'}^+(u') = 2$ and $d_{D'}(u) = d_{D'}(u') = 4$. On the other hand, there exists a vertex w in D' such that $w \notin \{u, u'\}$ and $N_{D'}^+(w) \cap \{u, u'\} \neq \phi$. We have $N_{D'}^+(w) = \{u, u'\}$ since D' contains no path P_4 and wt(v) cannot be an edge of G(D'); thus $d_{D'}(w) = 4$.

Since $d_{D'}(t(v)) = 4$, the vertices t(v), u, u' and w are in a block of D_4 which is neither complete nor a chordless odd cycle, which contradicts Theorem 2.11. This completes the proof of Theorem 2.1. \Box

An example which shows that the condition that each vertex has out-degree at least two in Theorem 2.1 is indispensable can be constructed from the 5-tournament T_5 with an edge (x, y) such that $x \notin V(T_5)$ and $y \in V(T_5)$.

If *H* contains a path P_4 , *x* cannot be an interior vertex of P_4 since d(x) = 1; furthermore it cannot be an end of P_4 since $d^-(x) = 0$. Thus $P_4 \subseteq T_5$ which contradicts Lemma 2.2.

We conclude this paper by asking the following question: Does there exist a 5-chromatic oriented graph which contains neither a 5-tournament nor P_4 ?

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